Asynchronous Shared Data Sources

Mart Lubbers (🖂) Haye Böhm Pieter Koopman Rinus Plasmeijer {mart,pieter,rinus}@cs.ru.nl haye.bohm@gmail.com

> LambdaDays 2021, 16–19 February 2021 TFP 2021, 17–19 February 2021





Data Sources

Data Sources



Shared Data Sources (SDSs)

Task Oriented Programming

Task Oriented Programming (TOP)

► Coordinate collaboration between people and machines to reach common goal.

- Declarative paradigm
- ▶ iTasks, mTasks, *TOP*
- Tasks are the basic building block
- Communication via Task Values and SDSs
- ► iTasks: event driven state transformers

SDSs in the iTask system

SDSs in iTasks

Programmer's interaction with shares

- ► get/watch
- ► set
- ► upd

. . .

. . .

User interaction with shares

- viewSharedInformation
- updateSharedInformation

System's reliance on shares

- events
- 7

:: Person = { name :: String , age :: Int , gender :: Gender } :: Gender = Male | Female | Other String

History

- ▶ Uniform Data Sources (Submitted for TFP 2012)
- Shared Data Sources (iTask system)
- ► Parametric Lenses (IFL 2014)
- Parametric Shared Data Sources (iTask system)

Limitations:

- Single thread
- Blocking
- Strict constraints
- Unsuitable for the real world



- Class based Shared Data Sources (MSc. Haye Böhm, this paper)
- Asynchronous Shared Data Sources (MSc. Haye Böhm, this paper)

Practical SDS use

Practical use: Slow network data

- MySQL asynchronous interface
- Web services

► OS specific waits (select, poll)



- ► Shares on a different machine
- Combinations of shares



Practical use: Asynchronous tasks

- Blocking tasks
- Tasks on different machines
- ► Tasks on different processes
- Communication via shares
- ► Tap directly into system shares



SDSs in general

Class based asynchronous SDS Operations

In general

Lenses	Parametrised Lenses
$\begin{array}{lll} get & \in X \to Y \\ put & \in Y \times X \to X \end{array}$	$\begin{array}{ll} get & \in \boldsymbol{\varPhi} \times X \to Y \\ put & \in \boldsymbol{\varPhi} \times Y \times X \to X \times (\boldsymbol{\varPhi} \to \textit{Bool}) \end{array}$

read

class read $v :: (v \ m \ p \ r \ w) \ p \rightarrow PViewT \ m \ r(ReadResult \ m \ p \ r \ w)$:: ReadResult $m \ p \ r \ w = Read \ r$ $\mid \exists sds: Reading (sds \ m \ p \ r \ w) \& read \ sds$

write

class write $v :: (v m p r w) p w \rightarrow PViewT m ()(WriteResult m p r w)$:: WriteResult m p r w = Written () $\mid \exists sds: Writing (sds m p r w) \& write sds$

Conclusion

Conclusion & Discussion

Conclusion

- Asynchronous Reads
- Asynchronous Writes
- Atomic updates*
- Used as we speak in the iTask system[†]

Discussion

- Complicated types[‡] for the compiler or we box them
- ▶ :: SDS m p r w = \exists sds: SDS (sds m p r w) & read sds & write sds

Questions?